

# Network Flows

UPCOPENCOURSEWARE number 34414

## Topic 1: Introduction

F.-Javier Heredia



**UNIVERSITAT POLITÈCNICA DE CATALUNYA**  
**BARCELONATECH**

**Departament d'Estadística  
i Investigació Operativa**



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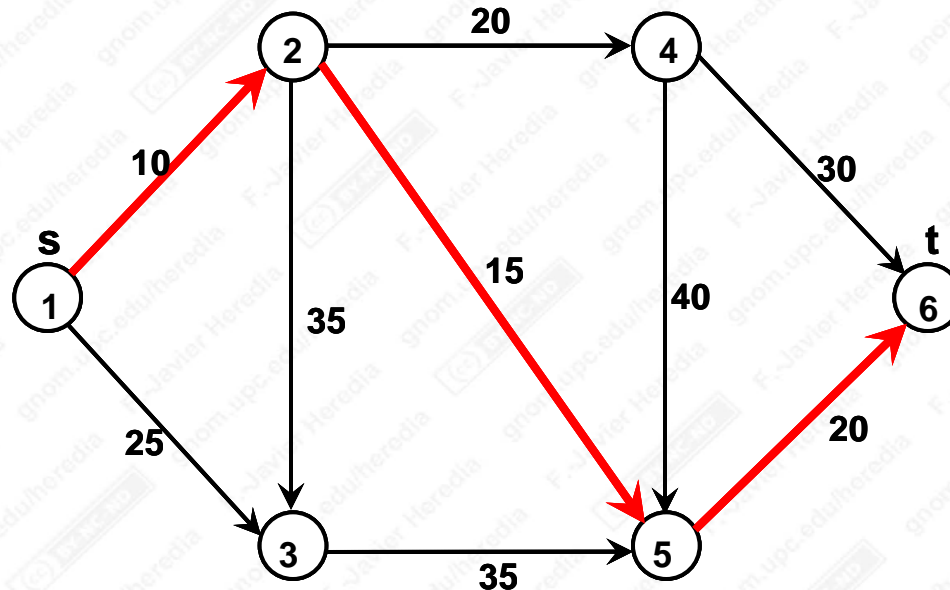
# Introduction

- **Definitions:**
  - Minimum shortest path problems.
  - Maximum flow problems.
  - Minimum cost flow problems.
- **Network flow problems (NF) as a class of linear programming problems (LP):**
  - Standard form of the minimum flow cost problem.
  - Nodes-arcs incidence matrix.
  - Formulation of the minimum shortest path problem.
  - Formulation of the maximum flow problem.
  - Formulation of the transportation and assignment problem.
  - Transformation to the standard form.
- Applications.



# Shortest Path Problems

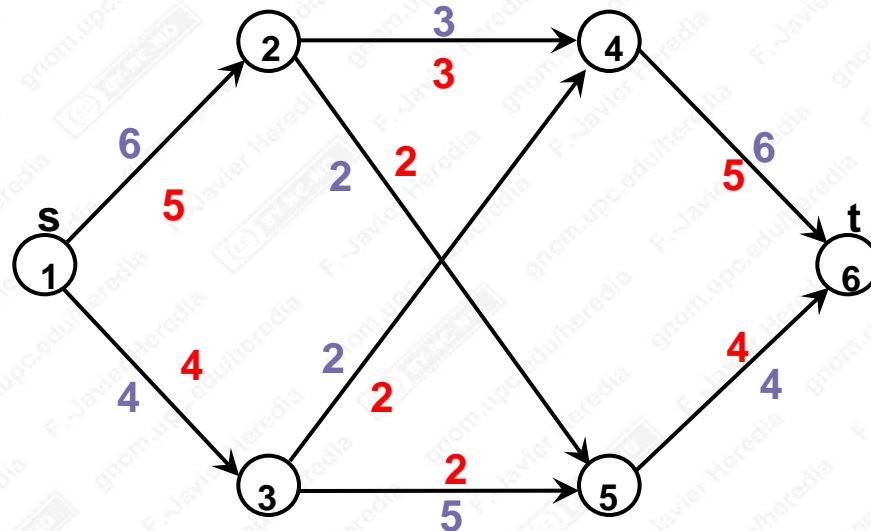
- To Identify the shortest path between a source node (“s”) and the sink node (“t”).



- Applications:
  - Minimum cost flow problem.
  - Minimum time flow problem.
  - Equipment replacement problems.

# Maximum Flow Problems

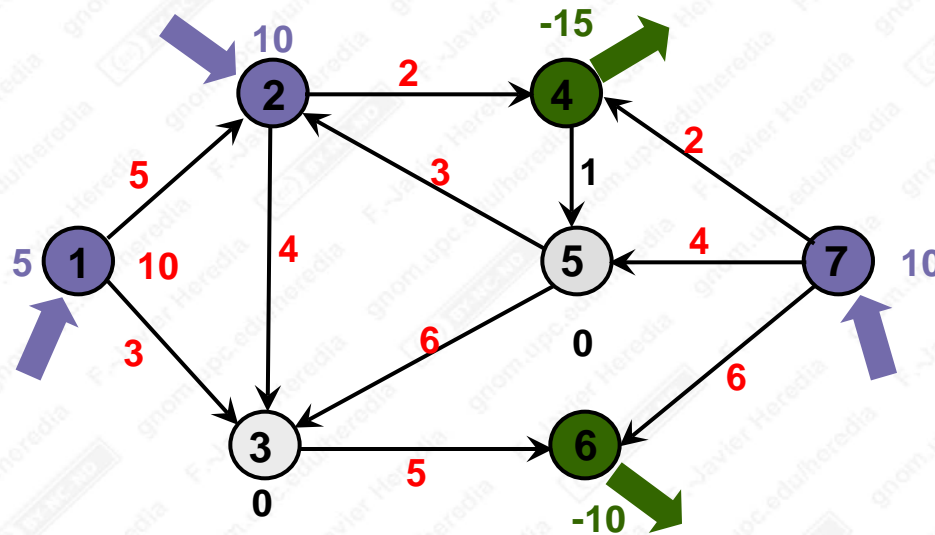
- To seek how to send the maximum possible amount of flow from a source node ("s") to sink node ("t") in a **capacitated network**.



- Applications: determining the maximum steady –state flow of:
  - Petroleum products in a pipeline network.
  - Cars in a road network.
  - Messages in a telecommunication network.
  - Electricity in an electrical network.

# Minimum Flow Cost Problems

- We wish to determine a least shipment cost of a single commodity through a network in order to satisfy consumption at **demand nodes** with the production of the **supply nodes**.

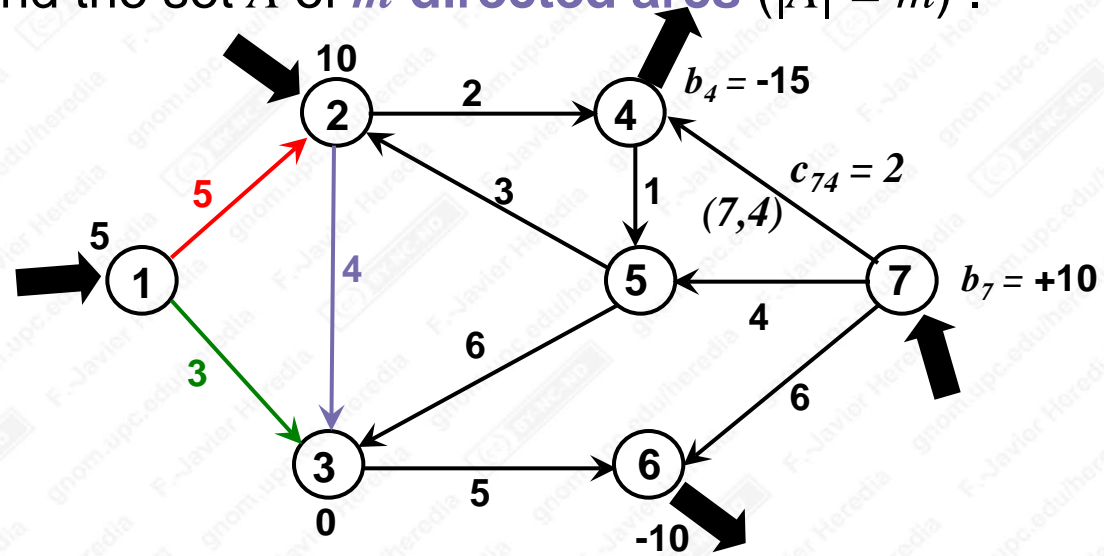


- Applications:
  - Logistics(warehouses to retailers).
  - Automobile routing in an urban traffic network.
  - Routing of calls through a telephone system.



# Standard form of the minimum cost flow problem (MCNFP) (I)

- Let  $G=(N,A)$  be the directed graph defined by the set  $N$  of  $n$  nodes ( $|N| = n$ ) and the set  $A$  of  $m$  directed arcs ( $|A| = m$ ).



$n=7 ; m=11 ; N=\{1,2,3,4,5,6,7\} ; A=\{ (1,2), (1,3), (2,3), \dots \}$

- Demand/supply vector:**  
 $b_j, j=1,2,\dots,n: b=[5,10,0,-15,0,5,10]'$
- Cost vector:**  $c_{ij}, (i,j) \in A : c=[5,3,4,\dots]'$
- Flow:**  $x_{ij}, (i,j) \in A$  : amount of commodity to be sent between node  $i$  and node  $j$  through arc  $(i,j)$ .

# Standard form of the minimum cost flow problem (II)

- Mathematic formulation:

$$\left\{ \begin{array}{l} \min \\ \text{s.a.:} \end{array} \right. \quad z = \sum_{(i,j) \in A} c_{ij} x_{ij} \quad (1.1a)$$

$$\sum_{\{j:(i,j) \in A\}} x_{ij} - \sum_{\{j:(j,i) \in A\}} x_{ji} = b_i \quad \forall i \in N \quad (1.1b)$$

**Balance equations**

$$x_{ij} \geq 0 \quad \forall (i, j) \in A \quad (1.1c)$$

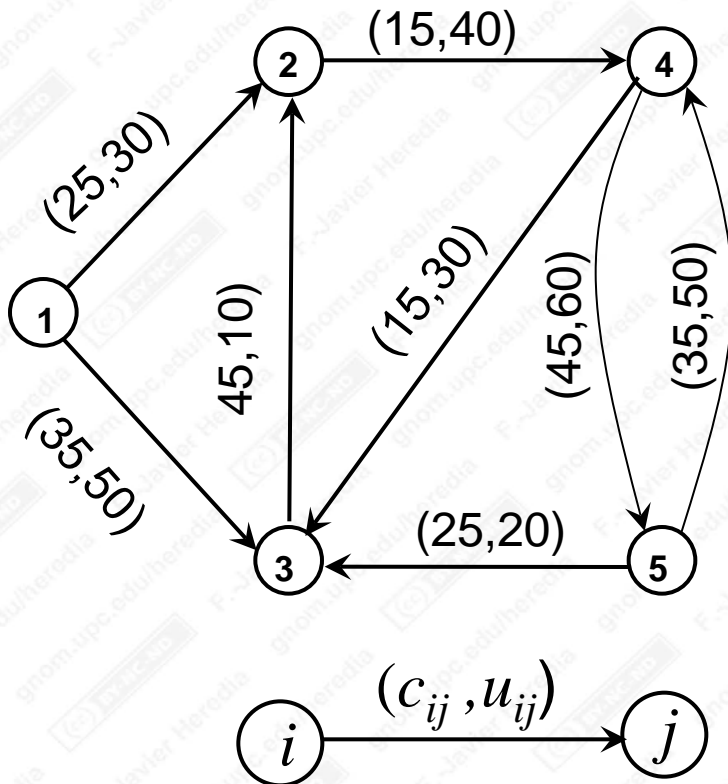
- We assume the network is **balanced**:  $\sum_{i=1}^n b_i = 0$

- Matrix notation:

$$\left\{ \begin{array}{l} \min \quad c'x \quad (1.2a) \\ \text{s.a.:} \quad Tx = b \quad (1.2b) \\ \quad \quad x \geq 0 \quad (1.2c) \end{array} \right.$$

# Node-arc incidence matrix

- Consider the following network:



- The associated node-arc incidence matrix  $T$  is:

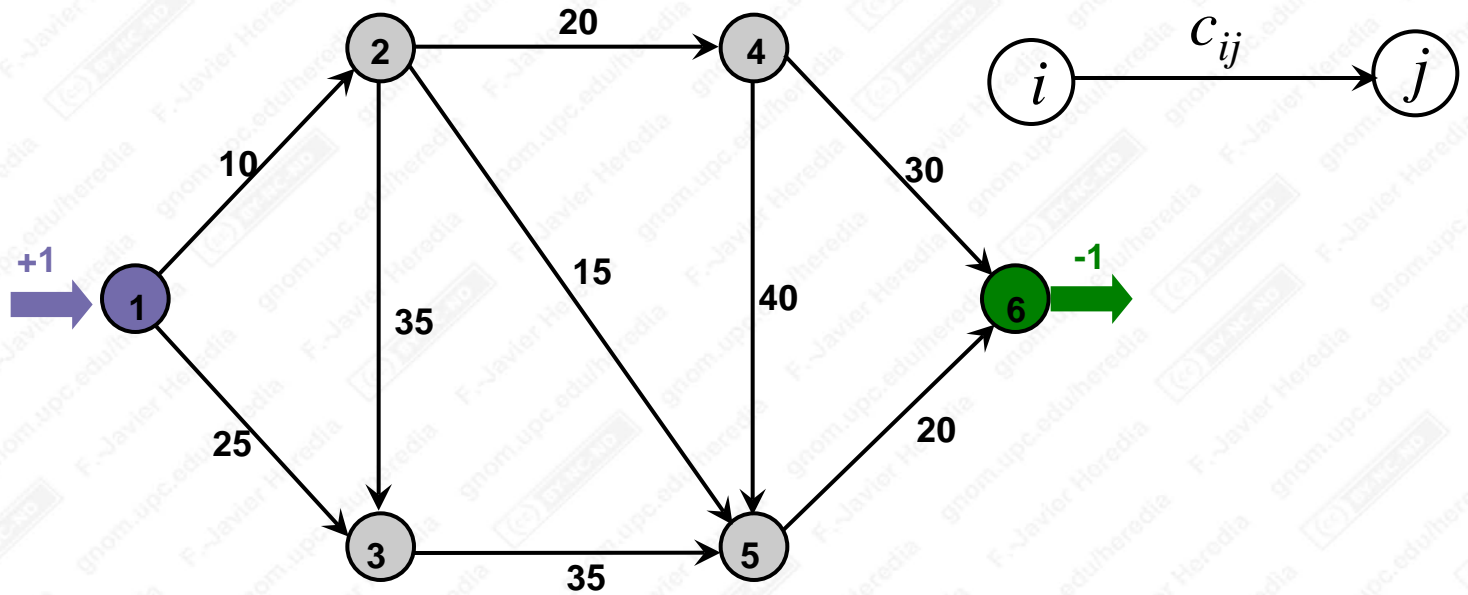
$$T = \begin{matrix} & \begin{matrix} (1,2) & (1,3) & (2,4) & (3,2) & (4,3) & (4,5) & (5,3) & (5,4) \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \\ 5 \end{matrix} & \begin{bmatrix} 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ -1 & 0 & 1 & -1 & 0 & 0 & 0 & 0 \\ 0 & -1 & 0 & 1 & -1 & 0 & -1 & 0 \\ 0 & 0 & -1 & 0 & 1 & 1 & 0 & -1 \\ 0 & 0 & 0 & 0 & 0 & -1 & 1 & 1 \end{bmatrix} \end{matrix}$$

- Properties:

- $2m$  non-zero elements among  $nm$ .
- Only 2 elements  $\neq 0$  per column.
- Every non-zero elements is +1 or -1.
- $\text{Rank}(T) = n-1$ .

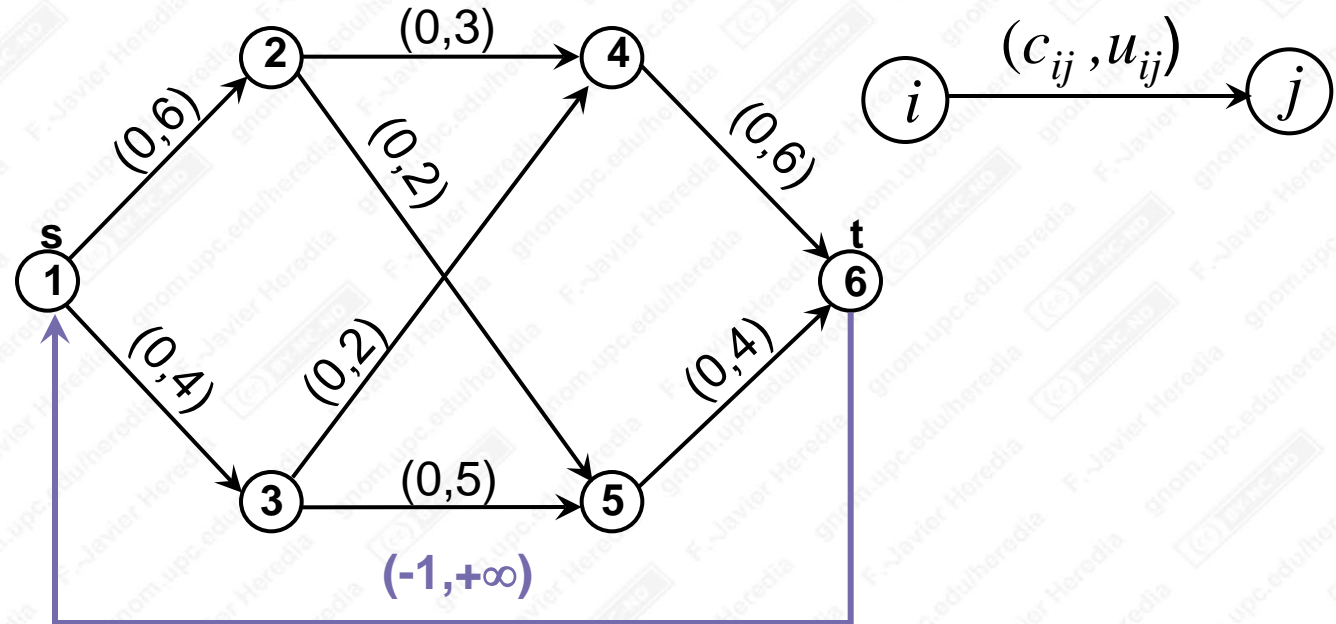


# Shortest path problem formulated as a MCFP



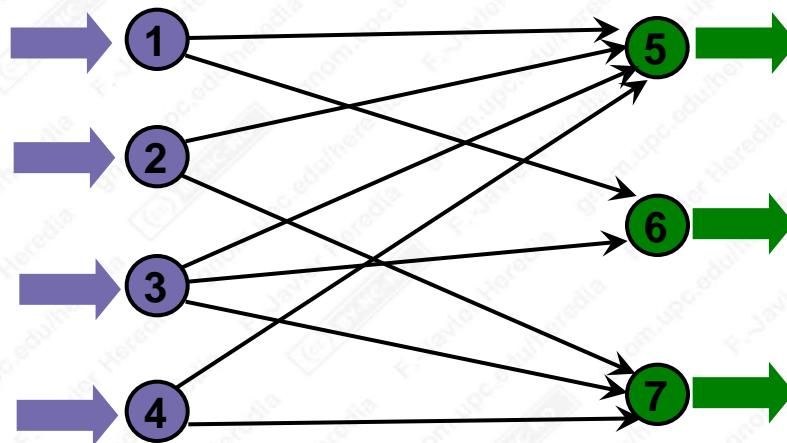
- $b_1 = +1 ; b_6 = -1 ; b_j = 0 \forall j \neq 1,6$

# Maximum flow problem formulated as a MCNFP



- Artificial arc  $x_{ts}$  with  $c_{ts} = -1$  and  $u_{ts} = +\infty$ .
- $c_{ij} = 0 \quad \forall (i, j) \neq (t, s)$

# Transportation problem formulated as a MCFP



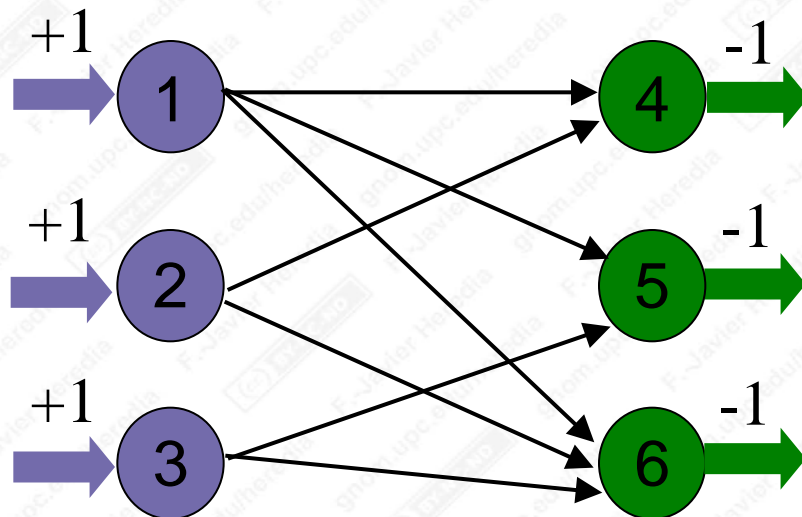
- Properties:

- $N = N_1 \cup N_2$ :

$N_1$ : production nodes;  $N_2$ : demand nodes.

- $\forall (i,j) \in A : i \in N_1 ; j \in N_2$

# Assignment problem formulation

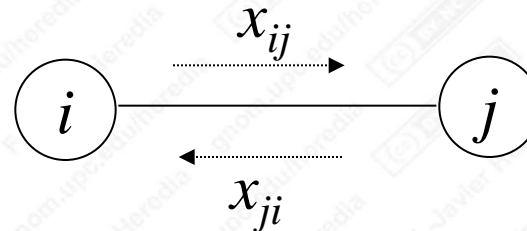


- Properties:

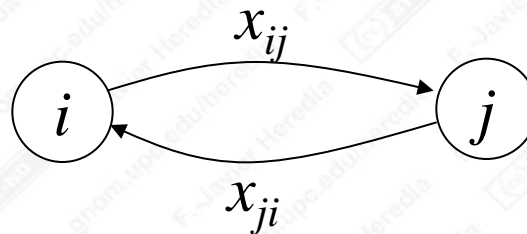
- $N = N_1 \cup N_2$  ;  $|N_1| = |N_2|$
- $A \subseteq N_1 \times N_2$ .
- $b_j = +1 \quad \forall j \in N_1$  ;  $b_j = -1 \quad \forall j \in N_2$

# Transformation to the standard form (I)

- **Undirected arcs:**



- Contribution to the objective function:  $c_{ij} x_{ij} + c_{ij} x_{ji}$
- $c_{ij} \geq 0 \Rightarrow$  at the optimal solution  $x_{ij} > 0$  or  $x_{ji} > 0$
- Transformation: each directed arc  $\{i, j\}$  is replaced by two directed arcs  $(i, j)$  and  $(j, i)$  with cost  $c_{ij}$  :





# Transformation to the standard form (II)

- **Arcs with non-zero lower bounds:**

- $x_{ij} \geq l_{ij}$

- $x_{ij}$  is replaced by  $x'_{ij} + l_{ij}$  in the problem's formulation:

- ❖ Bounds:  $x_{ij} \geq l_{ij} \Rightarrow x'_{ij} + l_{ij} \geq l_{ij} ; x'_{ij} \geq 0$

- ❖ Balance equations:  $b_i \rightarrow b_i - l_{ij} \quad ; \quad b_j \rightarrow b_j + l_{ij}$

- ❖ Objective function:  $c_{ij} x_{ij} \rightarrow c_{ij} (x'_{ij} + l_{ij}) = c_{ij} x'_{ij} + \boxed{c_{ij} l_{ij}}$   
cte.

# Transformation to the standard form (III)

- **Arcs with negative costs:**

- Let be  $x_{ij}$  with  $c_{ij} < 0$ .
- Let  $u_{ij}$  a trivial upper bound of the arc  $(i,j)$ .
- $x_{ij}$  is replaced by  $x'_{ij} = u_{ij} - x_{ij}$ .

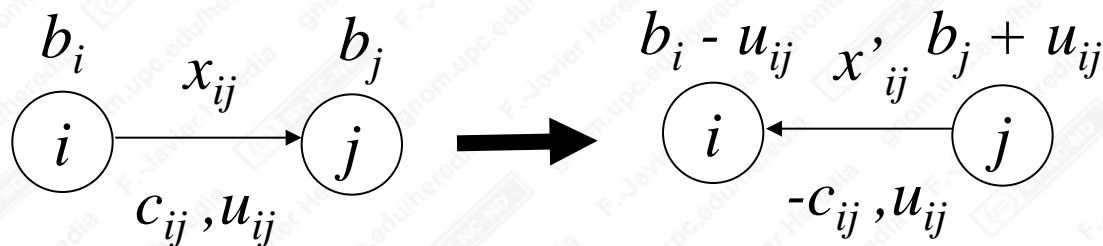
- ❖  $0 \leq x'_{ij} \leq u_{ij}$

- ❖ Objective function:

$$z = \dots + c_{ij} x_{ij} + \dots \rightarrow z' = \dots + \boxed{c_{ij} u_{ij}} \boxed{-c_{ij}} x'_{ij} + \dots$$

cst > 0

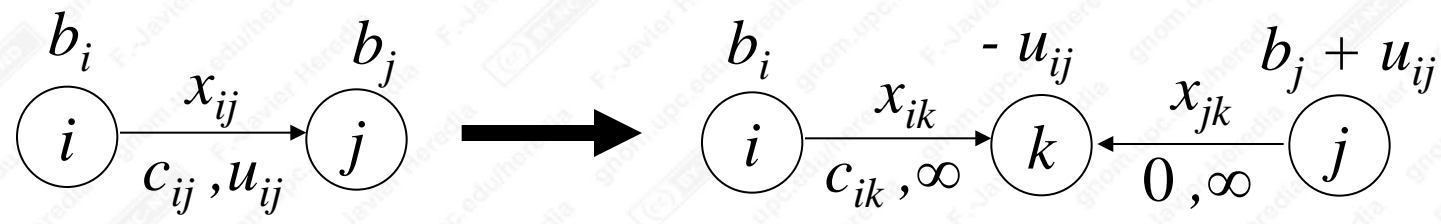
- ❖ Balance equations:



# Transformation to the standard form (IV)

- **Arcs with capacity:**

- Let  $0 \leq x_{ij} \leq u_{ij}$ .
- Network transformation:



❖ The network is equivalent to the original:

$$x_{ij} \equiv x_{ik} ; x_{ik} + x_{jk} = u_{ij} \implies x_{ik} = x_{ij} \leq u_{ij}$$

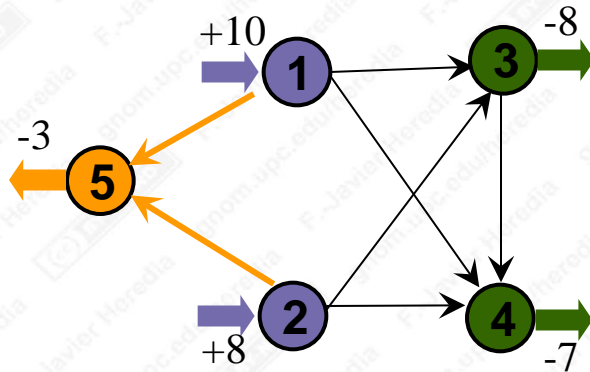
❖ The objective function is not modified:  $c_{ik} \equiv c_{ij}$

# Transformation to the standard form (V)

- **Unbalanced network:**

- **Excess of production**

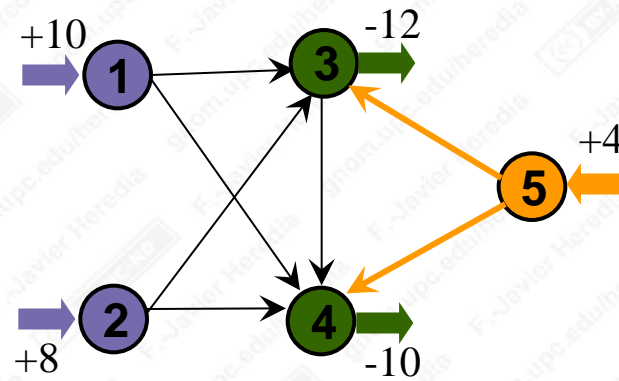
$$\left( \sum_{j=1}^n b_j > 0 \right) :$$



A **dummy demand node**  $n+1$  is added linking all production nodes through **uncapacitated - null cost arcs**.

- **Excess of demand**

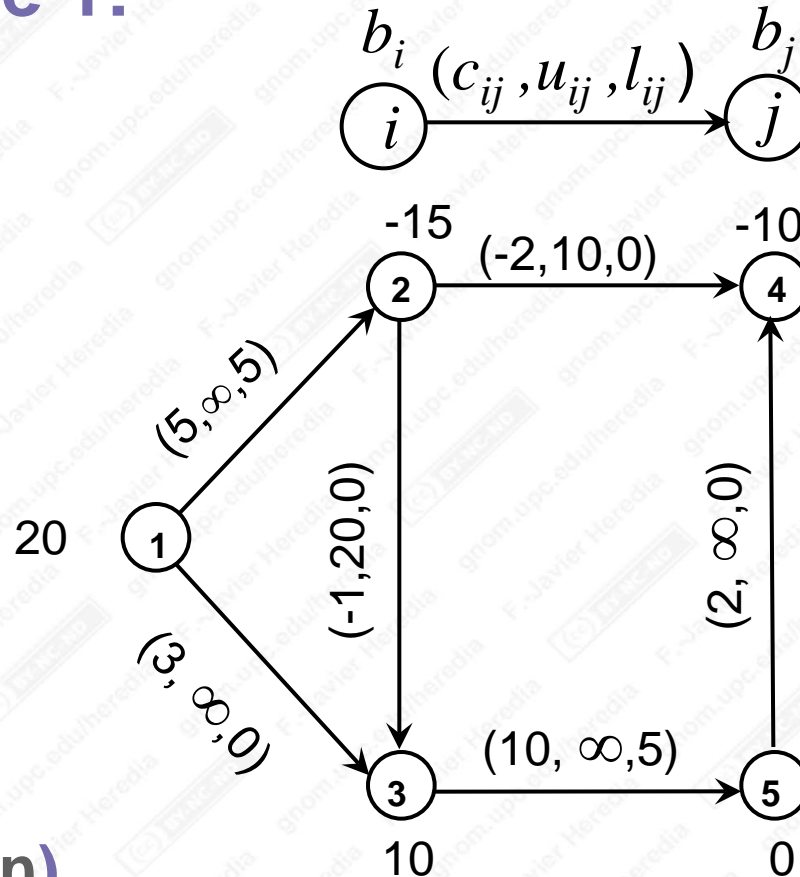
$$\left( \sum_{j=1}^n b_j < 0 \right) :$$



A **dummy supply node**  $n+1$  is added, linking all demand nodes through **uncapacitated - null cost arcs**.

# Transformation to the standard form (VI)

- Example 1:**



(Solution)





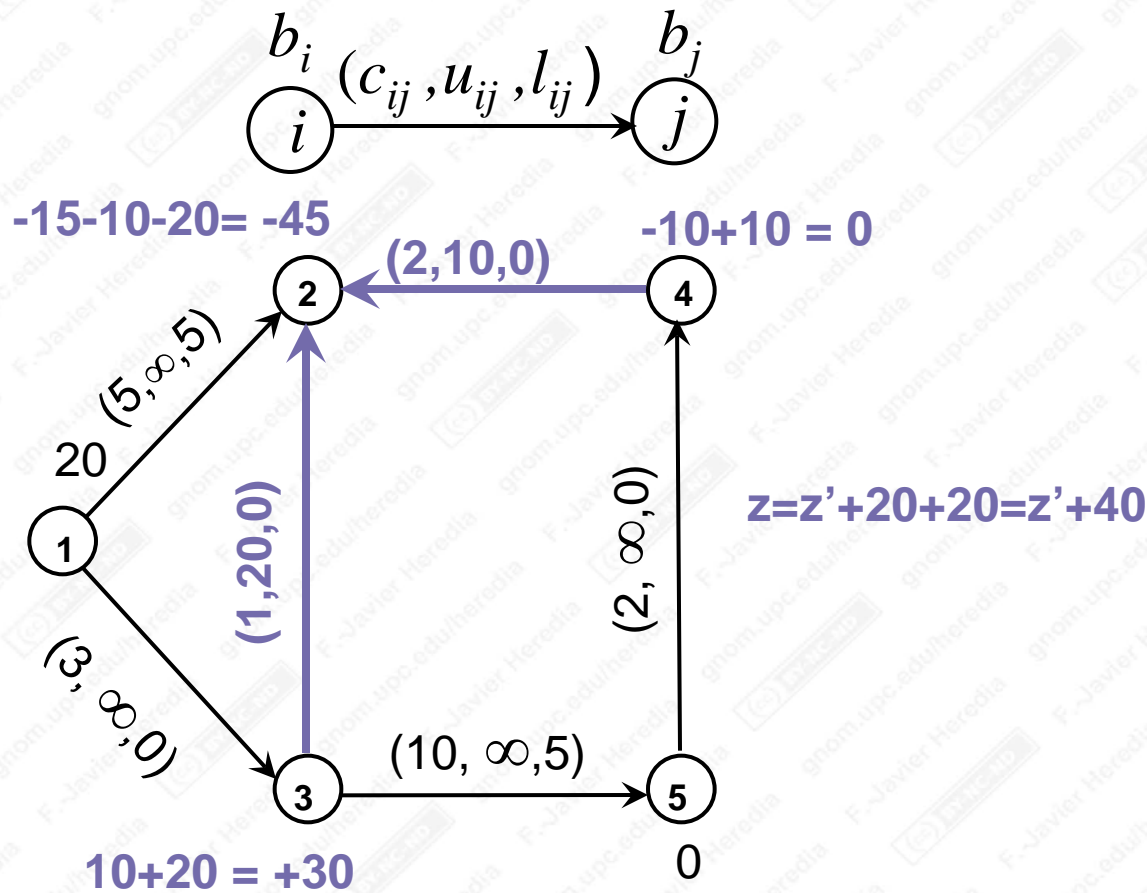
# Applications

- Solve the applications examples from book chapter 1.3 of AMO (1 week):
  - Problem description.
  - Network formulation and objective function associated to the problem.
  - Problem classification.
- Problems assignment:

# Transformation to the standard form

## Example (1/4)

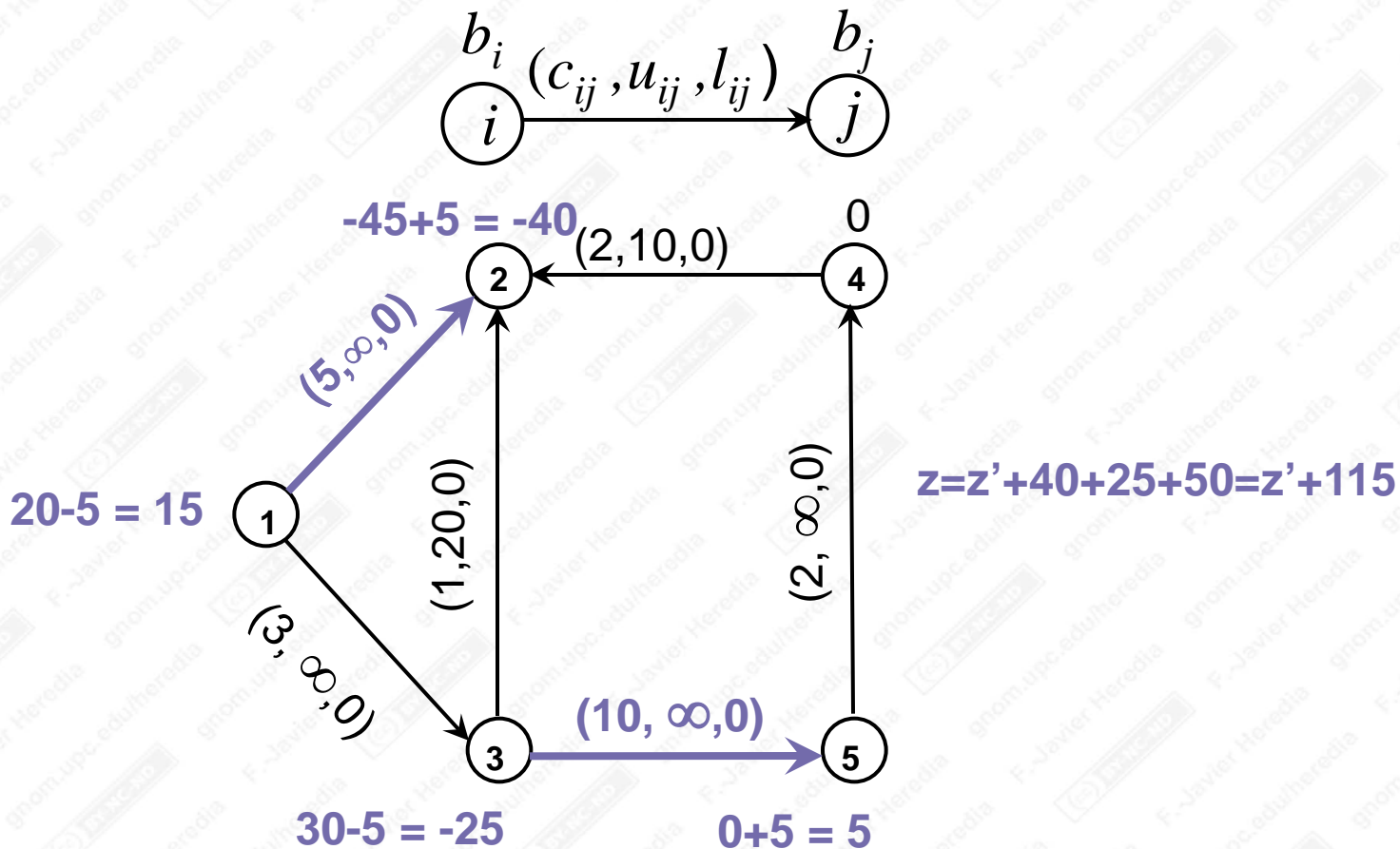
### 1. Negative cost elimination



# Transformation to the standard form

## Example (2/4)

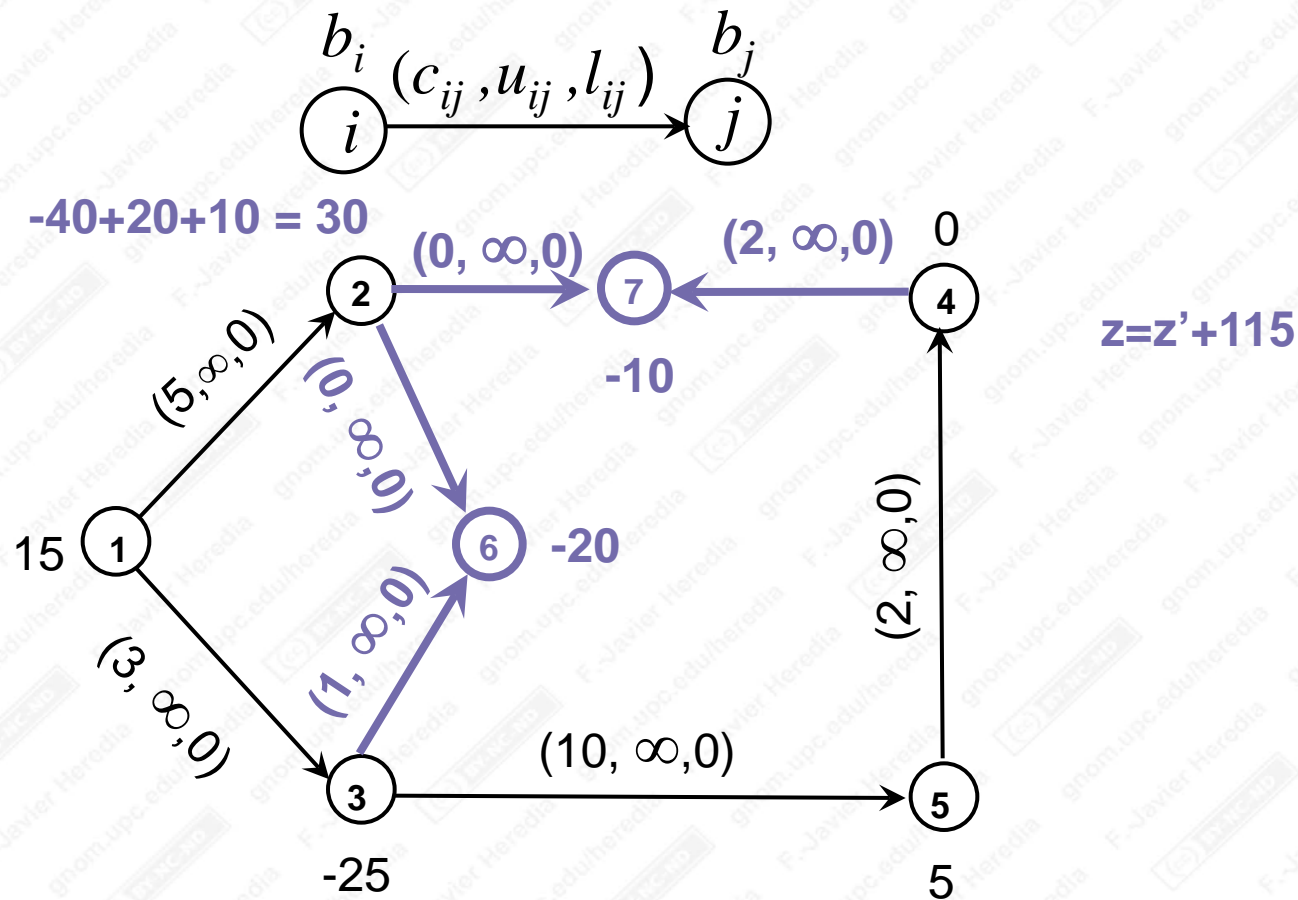
### 2. Lower bounds elimination $\neq 0$



# Transformation to the standard form

## Example (3/4)

### 3. Capacity elimination



# Transformation to the standard form

## Example (4/4)

### 4. Unbalance network

